MetaSync

File Synchronization Across Multiple Untrusted Storage Services

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File sync services are popular



400M of Dropbox users reached in June 2015

Many sync service providers











Can we rely on any single service?

Cloud Storage Often Results in Data Loss

All The Different Ways That 'iCloud' Naked Celebrity Photo Leak Might Have Happened







More

Shutting down Ubuntu One file services

U1, UBUNTU ONE

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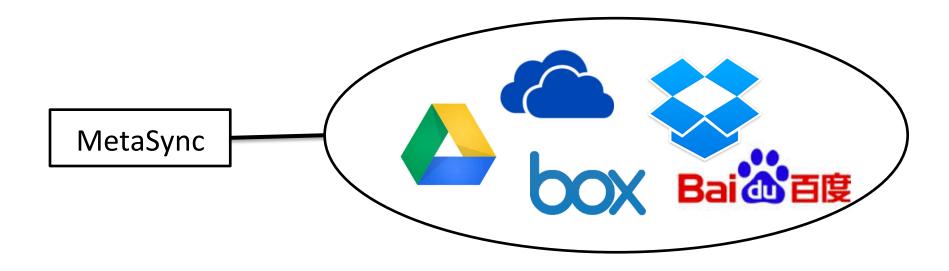
▲ Dropbox confirms that a bug within Selective Sync may have caused data loss (githubusercontent.com)

128 points by ghuntley 6 days ago | comments

Existing Approaches

- Encrypt files to prevent modification
 - Boxcryptor

- Rewrite file sync service to reduce trust
 - SUNDR (Li et al., 04), DEPOT (Mahajan et al., 10)



MetaSync:

Higher availability, greater capacity, higher performance Stronger confidentiality & integrity

Can we build a <u>better</u> file synchronization system across multiple existing services?

Goals

- Higher availability
- Stronger confidentiality & integrity
- Greater capacity and higher performance

- No service-service, client-client communication
- No additional server
- Open source software

Overview

- Motivation & Goals
- MetaSync Design
- Implementation
- Evaluation
- Conclusion

Key Challenges

 Maintain a globally consistent view of the synchronized files across multiple clients

 Using only the service providers' unmodified APIs without any centralized server

• Even in the *presence of service failure*

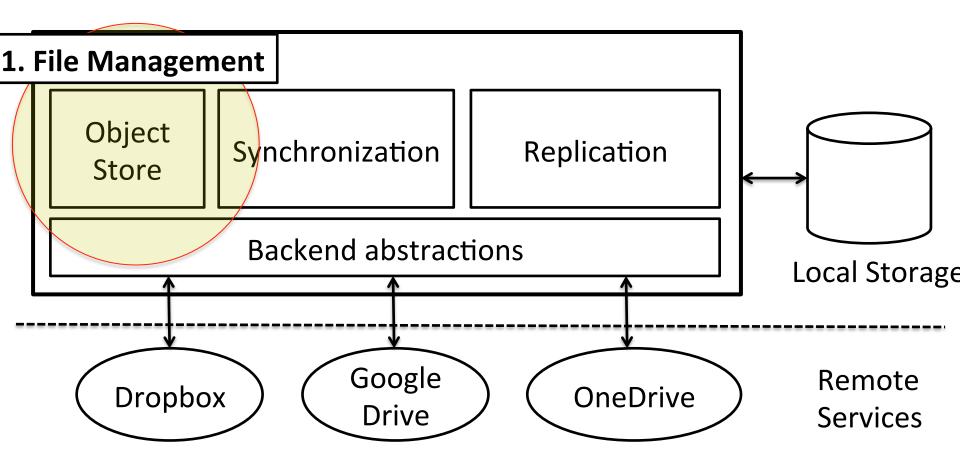
Design Choices

- How to manage files?
 - Content-based addressing & hash tree
- How to update consistently with unmodified APIs?
 - Client-based Paxos (pPaxos)
- How to spread files?
 - Stable deterministic mapping
- How to protect files?
 - Encryption from clients
- How to make it extensible?
 - Common abstractions

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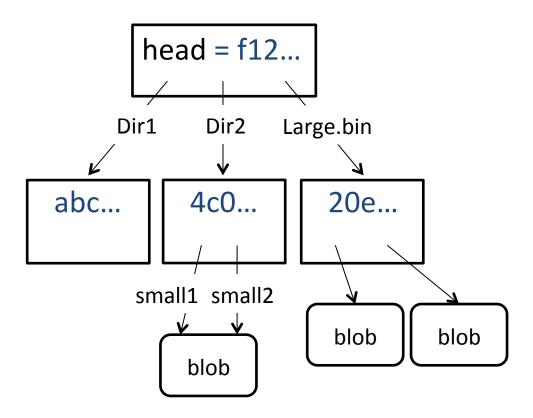
Overview of the Design



Object Store

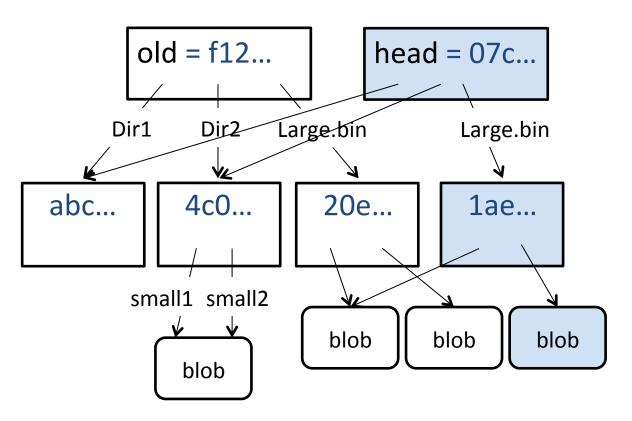
- Similar data structure with version control systems (e.g., git)
- Content-based addressing
 - File name = hash of the contents
 - De-duplication
 - Simple integrity checks
- Directories form a hash tree
 - Independent & concurrent updates

Object Store



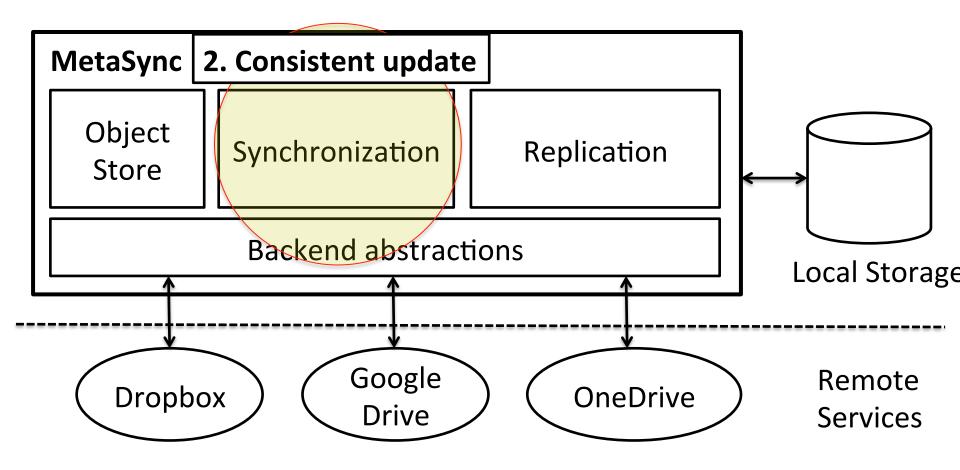
- Files are chunked or grouped into blobs
- The root hash = f12... uniquely identifies a snapshot

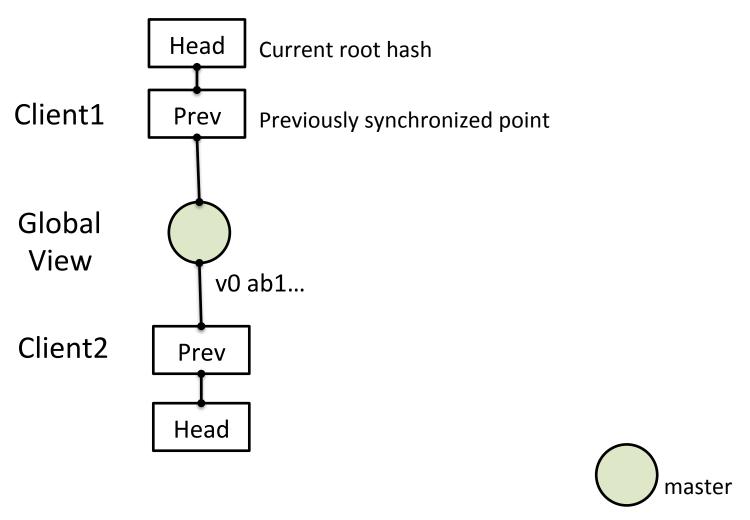
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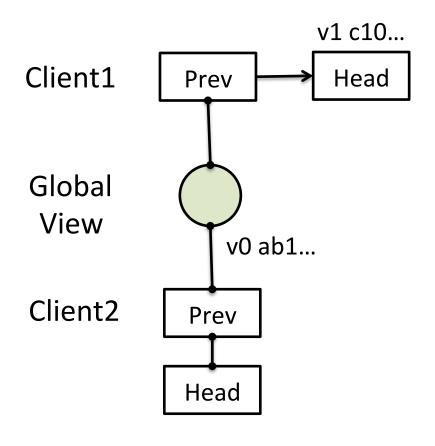


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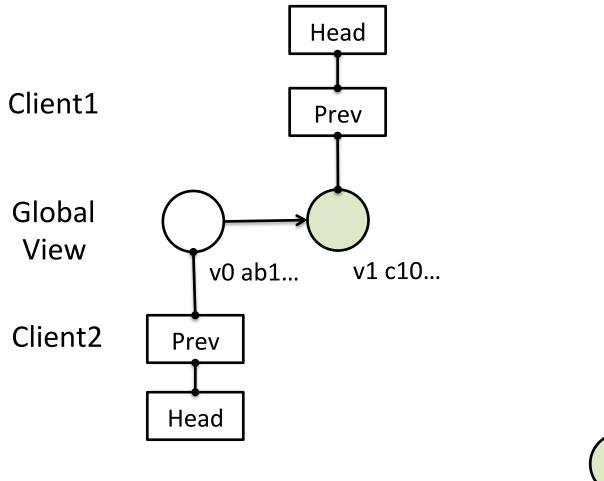
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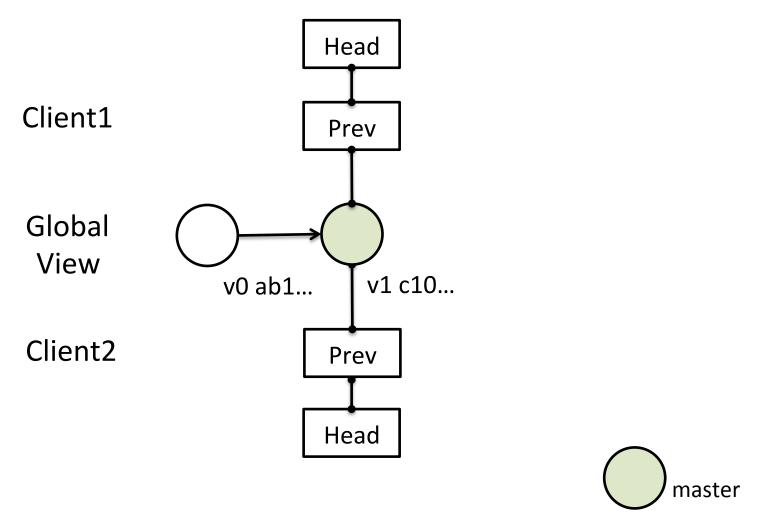


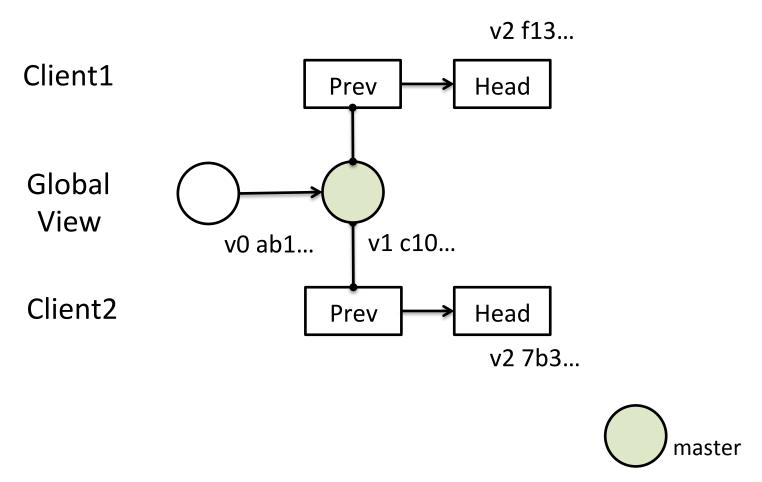


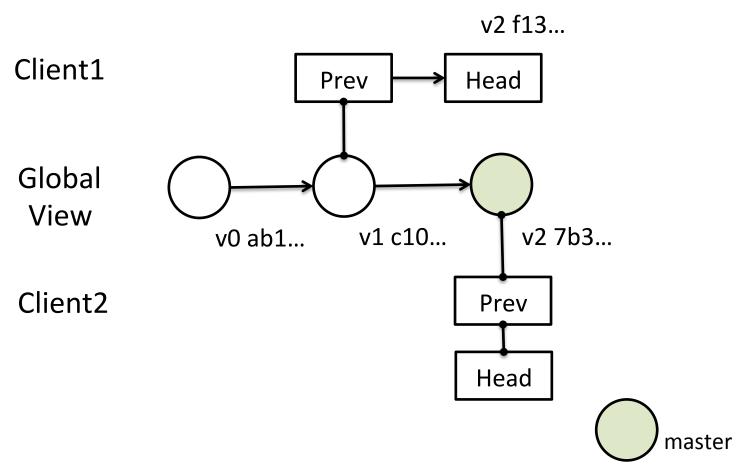






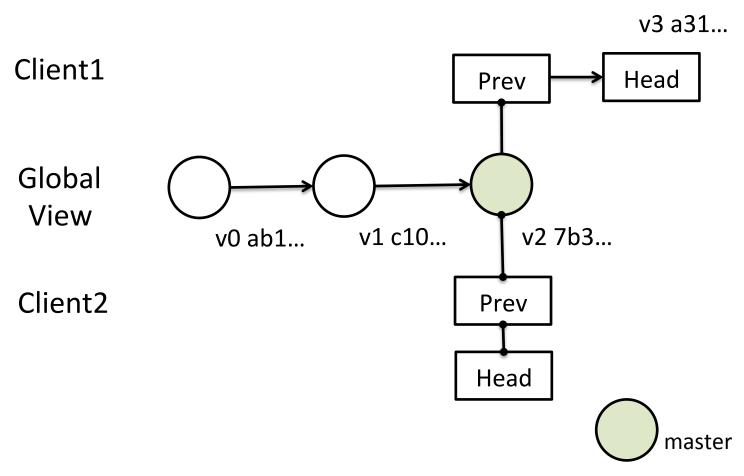




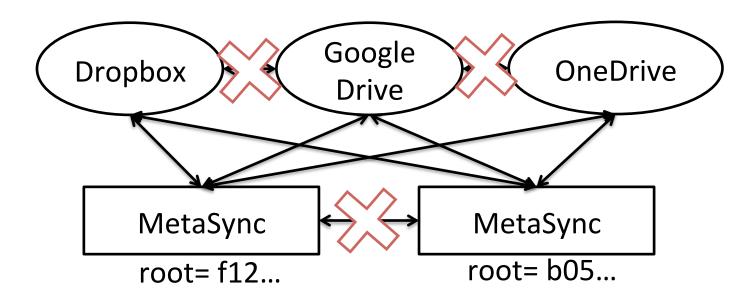


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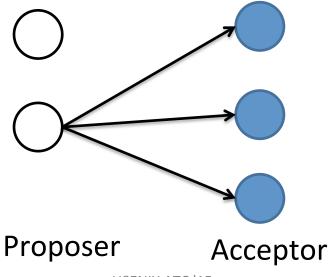
Consistent Update of Global View



 Need to handle concurrent updates, unavailable services based on existing APIs

Paxos

- Multi-round non-blocking consensus algorithm
 - Safe regardless of failures
 - Progress if majority is alive

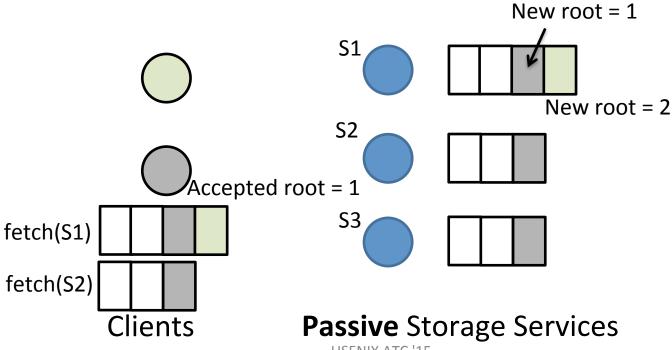


Metasync: Simulate Paxos

- Use an append-only list to log Paxos messages
 - Client sends normal Paxos messages
 - Upon arrival of message, service appends it into a list
 - Client can fetch a list of the ordered messages
- Each service provider has APIs to build appendonly list
 - Google Drive, OneDrive, Box: Comments on a file
 - Dropbox: Revision list of a file
 - Baidu: Files in a directory

Metasync: Passive Paxos (pPaxos)

- Backend services work as passive acceptor
- Acceptor decisions are delegated to clients

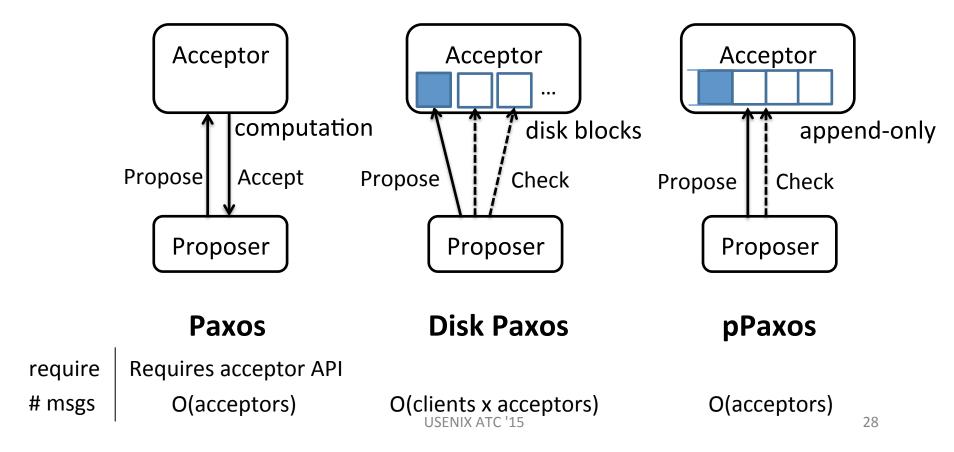


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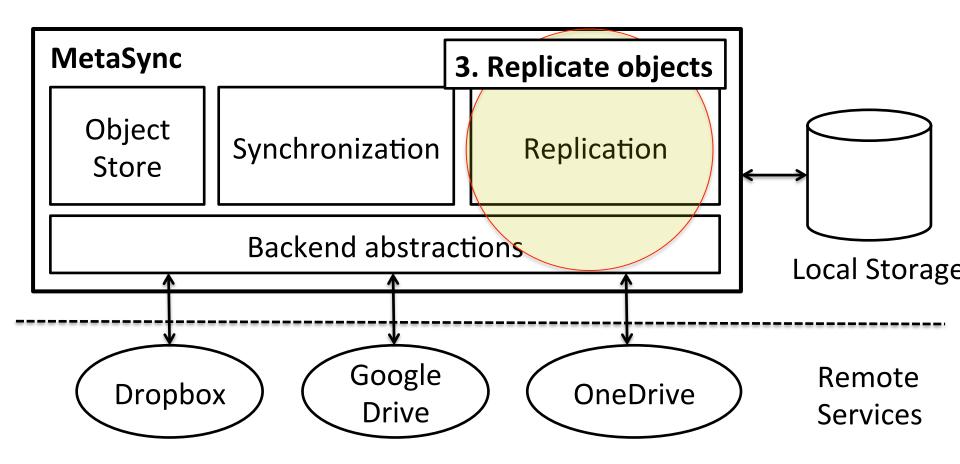
27

Paxos vs. Disk Paxos vs. pPaxos

Disk Paxos: maintains a block per client
 Gafni & Lamport '02



Overview of the Design



Stable Deterministic Mapping

- MetaSync replicates objects R times across S storage providers (R<S)
- Requirements
 - Share minimal information among services/clients
 - Support variation in storage size
 - Minimize realignment upon configuration changes
- Deterministic mapping

$$map: H \to \{s: |s| = R, s \subset S\}$$

- E.g., map(7a1...) = Dropbox, Google

Implementation

- Prototyped with Python
 - ~8k lines of code
- Currently supports 5 backend services
 - Dropbox, Google Drive, OneDrive, Box.net, Baidu
- Two front-end clients
 - Command line client
 - Sync daemon

Evaluation

How is the end-to-end performance?

 What's the performance characteristics of pPaxos?

 How quickly does MetaSync reconfigure mappings?

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End-to-End Performance

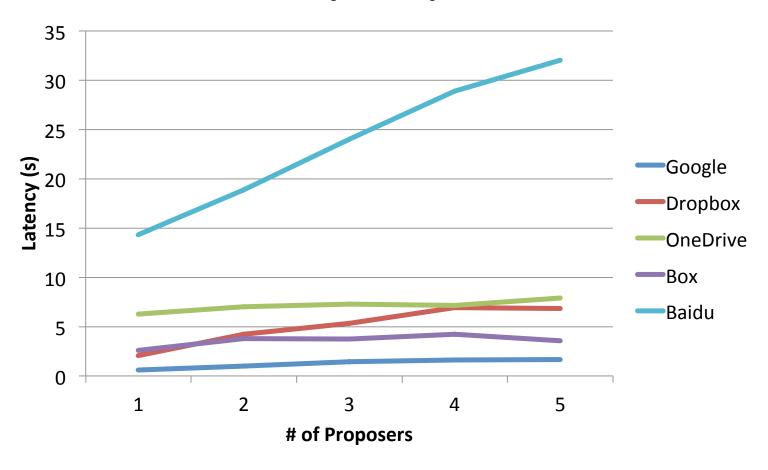
Synchronize the target between two computers

	Dropbox	Google	MetaSync
Linux Kernel 920 directories 15k files, 166MB	2h 45m	> 3hrs	12m 18s
Pictures 50 files, 193MB	415s	143s	112s
			(S = 4, R = 2)

Performance gains are from:

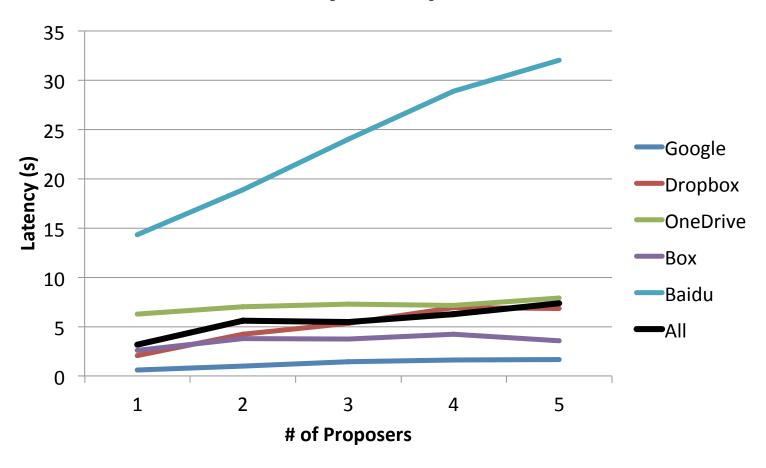
- Parallel upload/download with multiple providers
- Combined small files into a blob

Latency of pPaxos



Latency is not degraded with increasing concurrent proposers or adding slow backend storage service

Latency of pPaxos



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Conclusion

- MetaSync provides a secure, reliable, and performant files sync service on top of popular cloud providers
 - To achieve a consistent update, we devise a new client-based Paxos
 - To minimize redistribution, we present a stable deterministic mapping
- Source code is available:
 - http://uwnetworkslab.github.io/metasync/